Chapter 14: Planning with relational actions

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- Agents reason in time.
- Agents reason about time.
- Time passes as an agent acts and reasons.
- Given a goal, it is useful for an agent to think about what it will do in the future to determine what it will do now.
- Relational representations allow to reason about actions even before the individual objects of the domain become known to the agent.

Time can be modeled in a number of ways:

Discrete time Time is modeled as jumping from one time point to another.

Continuous time Time is modeled as being dense.

Event-based time Time steps don't have to be uniform; time steps can be between interesting events.

State space Instead of considering time explicitly, actions can map from one state to another.

You can model time in terms of points or intervals.

When modeling relations, you distinguish two basic types:

- Static relations are those relations whose value does not depend on time.
- Dynamic relations are relations whose truth values depends on time. Either
 - derived relations whose definition can be derived from other relations for each time,
 - primitive relations whose truth value can be determined by considering previous times.

The Delivery Robot World



- Individuals: rooms, doors, keys, parcels, and the robot.
- Actions:

move from room to room pick up and put down keys and packages unlock doors (with the appropriate keys)

 Relations: represent the robot's position the position of packages and keys and locked doors what the robot is holding

- at(Obj, Loc) is true in a situation if object Obj is at location *Loc* in the situation.
- carrying(Ag, Obj) is true in a situation if agent Ag is carrying Obj in that situation.
- sitting_at(Obj, Loc) is true in a situation if object Obj is sitting on the ground (not being carried) at location Loc in the situation.
- unlocked(Door) is true in a situation if door Door is unlocked in the situation.
- *autonomous*(*Ag*) is true if agent *Ag* can move autonomously. This is static.

- *opens(Key, Door)* is true if key *Key* opens door *Door*. This is static.
- *adjacent*(*Pos*₁, *Pos*₂) is true if position *Pos*₁ is adjacent to position *Pos*₂ so that the robot can move from *Pos*₁ to *Pos*₂ in one step.
- between(Door, Pos₁, Pos₂) is true if Door is between position Pos₁ and position Pos₂. If the door is unlocked, the two positions are adjacent.

- *move*(*Ag*, *From*, *To*) agent *Ag* moves from location *From* to adjacent location *To*. The agent must be sitting at location *From*.
- *pickup(Ag, Obj)* agent *Ag* picks up *Obj*. The agent must be at the location that *Obj* is sitting.
- *putdown(Ag, Obj)* the agent *Ag* puts down *Obj*. It must be holding *Obj*.
- *unlock*(*Ag*, *Door*) agent *Ag* unlocks *Door*. It must be outside the door and carrying the key to the door.

sitting_at(rob, o109).
sitting_at(parcel, storage).
sitting_at(k1, mail).

between(door1, o103, lab2).
opens(k1, door1).
autonomous(rob).

```
at(Obj, Pos) \leftarrow sitting\_at(Obj, Pos).

at(Obj, Pos) \leftarrow carrying(Ag, Obj) \land at(Ag, Pos).

adjacent(o109, o103).

adjacent(o103, o109).

\cdots

adjacent(lab2, o109).

adjacent(P_1, P_2) \leftarrow

between(Door, P_1, P_2) \land
```

unlocked(Door).

- Often we want to refer to individuals in terms of components.
- Examples: 4:55 p.m. English sentences. A classlist.
- We extend the notion of term. So that a term can be $f(t_1, ..., t_n)$ where f is a function symbol and the t_i are terms.
- In an interpretation and with a variable assignment, term $f(t_1, \ldots, t_n)$ denotes an individual in the domain.
- Function symbols can be recursively embedded:
 - One function symbol and one constant can refer to infinitely many individuals.
 - Function symbols can be used to describe sequential structures.

- A list is an ordered sequence of elements.
- Let's use the constant *nil* to denote the empty list, and the function *cons(H,T)* to denote the list with first element *H* and rest-of-list *T*. These are not built-in.
- The list containing sue, kim and randy is

cons(sue, cons(kim, cons(randy, nil)))

append(X, Y, Z) is true if list Z contains the elements of X followed by the elements of Y append(nil, Z, Z).
 append(cons(A, X), Y, cons(A, Z)) ← append(X, Y, Z).

Situation Calculus

- State-based representation where the states are denoted by terms.
- A situation is a term that denotes a state.
- There are two ways to refer to states:
 - *init* denotes the initial state
 - do(A, S) denotes the state resulting from doing action A in state S, if it is possible to do A in S.
- Time is reified by means of situations which can be entered.
 - Individuals represent points in time.
- A situation encodes how to get to the state it denotes.
 - A state may be represented by multiple situations.
 - A state may be represented by no situations if it is unreachable.
 - A situation may represent no states, if an action was not possible.

init

- do(move(rob, o109, o103), init)
- do(move(rob, o103, mail), do(move(rob, o109, o103), init)).
- do(pickup(rob, k1), do(move(rob, o103, mail), do(move(rob, o109, o103), init))).

- Add an extra term to each dynamic predicate indicating the situation.
- Example Atoms:

at(rob, o109, init)

at(rob, o103, do(move(rob, o109, o103), init))
at(k1, mail, do(move(rob, o109, o103), init))

- You specify what is true in the initial state using axioms with *init* as the situation parameter.
- Primitive relations are axiomatized by specifying what is true in situation do(A, S) in terms of what holds in situation S.
- Derived relations are defined using clauses with a free variable in the situation argument.
- Static relations are defined without reference to the situation.

. . .

sitting_at(rob, o109, init).
sitting_at(parcel, storage, init).
sitting_at(k1, mail, init).

```
adjacent(P_1, P_2, S) \leftarrow
between(Door, P_1, P_2) \land
unlocked(Door, S).
adjacent(lab2, o109, S).
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 $\begin{array}{l} \hline poss(A,S) & \text{is true if action } A \text{ is possible in situation } S.\\ \hline poss(putdown(Ag, Obj), S) \leftarrow\\ & carrying(Ag, Obj, S).\\ \hline poss(move(Ag, Pos_1, Pos_2), S) \leftarrow\\ & autonomous(Ag) \wedge\\ & adjacent(Pos_1, Pos_2, S) \wedge\\ & sitting_at(Ag, Pos_1, S). \end{array}$

Example: Unlocking the door makes the door unlocked: $unlocked(Door, do(unlock(Ag, Door), S)) \leftarrow poss(unlock(Ag, Door), S).$

Frame Axiom: No actions lock the door:

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unlocked(Door, do(A, S)) \leftarrow
unlocked(Door, S) \land
poss(A, S).
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Picking up an object causes it to be carried:

 $carrying(Ag, Obj, do(pickup(Ag, Obj), S)) \leftarrow poss(pickup(Ag, Obj), S).$

Frame Axiom: The object is being carried if it was being carried before unless the action was to put down the object:

 $carrying(Ag, Obj, do(A, S)) \leftarrow$ $carrying(Ag, Obj, S) \land$ $poss(A, S) \land$ $A \neq putdown(Ag, Obj).$

An object is sitting at a location if:

• it moved to that location:

 $sitting_at(Obj, Pos, do(move(Obj, Pos_0, Pos), S)) \leftarrow poss(move(Obj, Pos_0, Pos).$

• it was put down at that location:

sitting_at(Obj, Pos, do(putdown(Ag, Obj), S)) \leftarrow poss(putdown(Ag, Obj), S) \land at(Ag, Pos, S).

 it was at that location before and didn't move and wasn't picked up. The only actions that undo *sitting_at* for object *Obj* is when *Obj* moves somewhere or when someone is picking up *Obj*.

$$\begin{array}{l} \textit{sitting_at(Obj, Pos, do(A, S))} \leftarrow \\ \textit{poss}(A, S) \land \\ \textit{sitting_at}(Obj, Pos, S) \land \\ \forall \textit{Pos}_1 \quad A \neq \textit{move}(Obj, \textit{Pos}, \textit{Pos}_1) \land \\ \forall \textit{Ag} \quad A \neq \textit{pickup}(\textit{Ag}, \textit{Obj}). \end{array}$$

$$\forall Ag \ A \neq pickup(Ag, Obj)$$

is equivalent to:

$$\sim \exists Ag \ A = pickup(Ag, Obj)$$

which can be implemented as

sitting_at(Obj, Pos, do(A, S))
$$\leftarrow \cdots \land \cdots \land \cdots \land \cdots \land$$

 \sim is_pickup_action(A, Obj).

with the clause:

 $is_pickup_action(A, Obj) \leftarrow A = pickup(Ag, Obj).$

which is equivalent to:

is_pickup_action(pickup(Ag, Obj), Obj).

Given

- an initial world description
- a description of available actions
- a goal
- a plan is a sequence of actions that will achieve the goal.

If you want a plan to achieve Rob holding the key k1 and being at o103, the query

? $carrying(rob, k1, S) \land at(rob, o103, S)$.

has an answer

S = do(move(rob, mail, o103), do(pickup(rob, k1), do(move(rob, o103, mail), do(move(rob, o109, o103), init)))).

The Delivery Robot World



- Idea: backward chain on the situation calculus rules.
- A complete search strategy (e.g., A* or iterative deepening) is guaranteed to find a solution.
- When there is a solution to the query with situation $S = do(A, S_1)$, action A is the last action in the plan.
- You can virtually always use a frame axiom so that the search space is largely unconstrained by the goal. Search space is enormous.

• Given a goal, you would like to consider only those actions that actually achieve it.

Example:

? carrying(rob, parcel, S) \land in(rob, lab2, S).

the last action needed is irrelevant to the left subgoal.

• So we need to combine the planning algorithms with the relational representations.